



Truncated Random Measures

Jonathan Huggins

MIT CSAIL and Dept. of EECS

with: T. Campbell, J. How, T. Broderick





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- Goal: integrate BNP priors into PPLs like Stan



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Bayesian nonparametrics:

achieves growing model size via infinite parameters

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New Hork

Eimes

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automate inference with probabilistic programming

[Gopalan 2014] [Teh 2006] [Huang 2014] [Michini 2015] [Lennox 2010] [Prunster 2014] [Yang 2015] [Yu 2012] [Ozaki 2008] [Kottas 2008]

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Wide variety of priors in

BNP with no finite

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Previously studied priors	

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 2 representation forms (7 reps total) that allow finite approximation of (normalized) completely random measures ((N)CRMs)

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Contributions:

- All BNP priors Priors with finite approx (new) Previously studied priors with finite approx (past work)
- 2 representation forms (7 reps total) that allow finite approximation of (normalized) completely random measures ((N)CRMs)
- Approximation error analysis

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- Approximation error analysis
- Computational complexity analysis (not in this talk)

	Finite Approximation	Approximation Error Bounds	Computational Complexity
DP	\checkmark	▶	▶
BP	\checkmark	V	✓
BPP	♥		
ΓΡ	\checkmark	\checkmark	V
(N)CRM			

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DP	[Sethuraman 94] [Roychowdhury 15]	▶ [Ishwaran 01]	♥
BP	[Teh 07] [Paisley 12] [Thibaux 07]	[Doshi-Velez 09][Paisley 12]	V
BPP	[Broderick 14]		
ΓΡ	[Bondesson 82] [Roychowdhury 15]	✓ [Roychowdhury 15]	V
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(N)CRM	[Brodering]	general th	eory
Truncation Roadmap

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Tractable models in BNP





Truncation Roadmap



























sports

topic

space



0.7 0.5 0.2

















How do we generate infinitely many trait/rate points (ψ , θ)?



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trait space

[Kingman 93]

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Captures a large class of useful priors in BNP

How do we generate infinitely many trait/rate points (ψ , θ)?



How do we pick a finite subset of the points?

[Kingman 93]





We pick a finite subset of atoms (ψ , θ) by:



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1) ordering the atoms (sequential representation)





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Series representation

function of a homogenous Poisson point process (4 versions)

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Superposition representation

infinite sum of homogenous CRMs, each with finite # of atoms (3 versions)

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function of a homogenous Poisson point process (4 versions) Superposition representation

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Theorem (H., Campbell, How, Broderick). Can generate (N)CRMs using all 7 sequential representations

Sequential representation comparison

Why so many representations?

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They're all useful in different circumstances

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	Series Reps				Superposition Reps		
	B-Rep	IL-Rep	R-Rep	T-Rep	DB-Rep	PL-Rep	SB-Rep
Error Bound Decay	(exp)	(exp)	√ / X	X	(exp)	(exp)	X
Ease of Analysis	X	XX	X	X	\checkmark	\checkmark	\checkmark
Generality	V	\checkmark	✓	✓	V	✓	\checkmark
Known # Atoms	\checkmark	\checkmark	X	X	X	X	X

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Step 1: compute $c := \lim_{\theta \to 0} \theta \nu(\theta) = \gamma \lambda$ **Step 2:** compute $f(\theta) := -c^{-1} \frac{d}{d\theta} [\theta \nu(\theta)] = \lambda e^{-\lambda \theta}$ **Step 3:** plug in! $\Theta = \sum_{k=1}^{\infty} V_k e^{-\Gamma_k} \delta_{\psi_k}, \quad V_k \stackrel{\text{iid}}{\sim} f, \quad \Gamma \sim \text{PoissonP}(c)$





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full infinite O truncated Θ_{κ}

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Compare the distribution of the data under full vs. truncated

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Cannot evaluate exactly, so we develop new upper bounds

Protobound mist

Leads to all the other truncation error bounds in this work



Protobound missik

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Previous Work

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BP	\checkmark	▶	√
BPP	▼		
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BPP	⋎		
ΓΡ	\checkmark	\checkmark	V
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(N)CRM	\checkmark	\checkmark	

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J. Huggins^{*}, T. Campbell^{*}, J. How, T. Broderick. **Truncated Random Measures.** Submitted, 2016. Available online: <u>https://arxiv.org/abs/1603.00861</u>